

A binary relation \circlearrowleft on set $A \subset A \times A$ $A =$ a choice set.

$(x, y) \in \circlearrowleft$ or $x \circlearrowleft y$ \longleftarrow Proposition forming; e.g. taller

e.g. arithmetic " $<$ " is a binary relation

" $+$ " is not a binary relation, $x + y$ is not a Proposition

Consistency Conditions

A Preference ordering \geq_A is a binary relation with properties

1. Complete. $\forall x, y \ x \geq_A y$ or $y \geq_A x$.

2. Transitive $\forall x, y, z$, if $x \geq_A y$ and $y \geq_A z \Rightarrow x \geq_A z$.

3. Independent of irrelevant alternatives

$$\forall x, y \in B, \ x \underset{B}{\geq} y \iff x \underset{A}{\geq} y.$$

x is weakly preferred to $y \equiv$ An individual chooses $x \in A$ such that $\forall y \in A, \ x \geq y$.

Completeness implies reflexivity.

Strong Preference is denoted by \succ .

Weak Preference is " " \succsim .

$x \sim y \equiv$ equivalent iff $x \succsim y$ & $y \succsim x$.

exclusion condition \equiv if $x \succ y$ then $\neg(y \succ x)$.

A Preference relation is consistent iff it meets ^{the} consistency conditions.

Theorem: A binary relation \geq on A of payoffs can be represented by a preference function $u: A \rightarrow \mathbb{R}$ iff \geq is consistent.

Exponential Discounting

Let x & y be two consumption streams

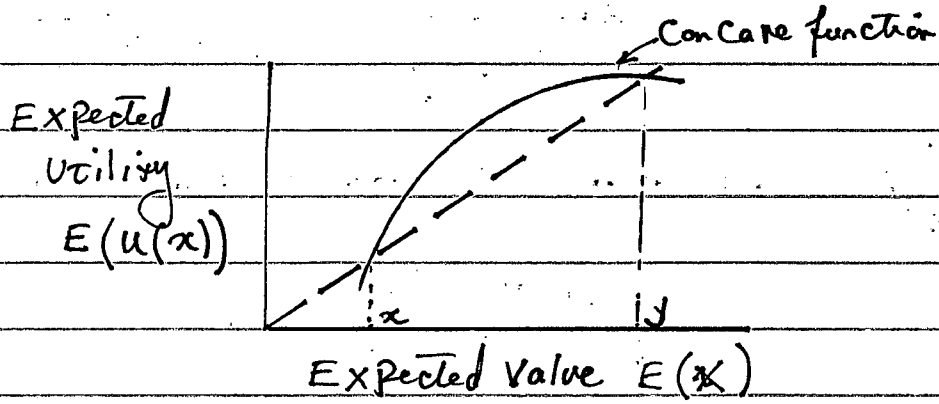
$$x = x_0, x_1, \dots \quad y = y_0, y_1, \dots$$

Let $\delta \in (0, 1)$ be a constant capturing a discount factor.

$$U(x_0, x_1, \dots) = \sum_{k=0}^{\infty} \delta^k u(x_k)$$

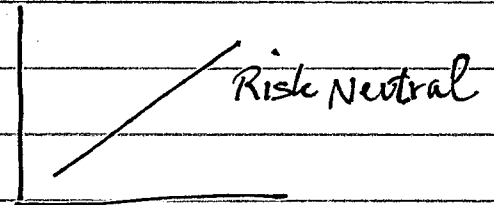
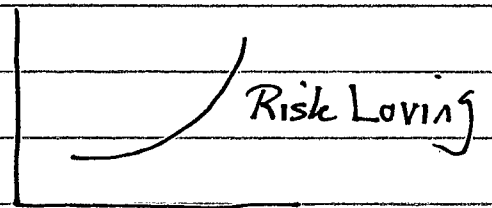
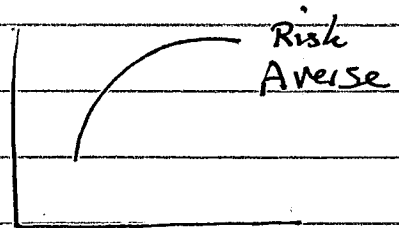
Hexmoor
SIUC

If $X = \mathbb{R}$, Payoffs can be considered to be money.



$\forall x, y \in X, p u(x) + (1-p) u(y)$

$$p u(x) + (1-p) u(y) < u(p x + (1-p) \cdot y)$$



A normal form game

		II		
		L	M	R
I	T	(1,1)	(0,2)	(2,1)
	M	(2,2)	(1,1)	(0,0)
	B	(1,0)	(0,0)	(-1,1)

GT
L2
Van Neuman/
Morgenstern
Utilities for
players I & 2.

Payoffs are Common Knowledge.

For I, $M > B$, For II, \nexists a dominant Strategy.

Stable outcome \equiv Strategy Profiles Where no player has
Incentive to deviate if he knows that the other players play the
Prescribed Strategies.

E.g.,

		L	R
I	T	(1,1) ^o	(0,0) [*]
	B	(0,0) [*]	(1,1) ^o

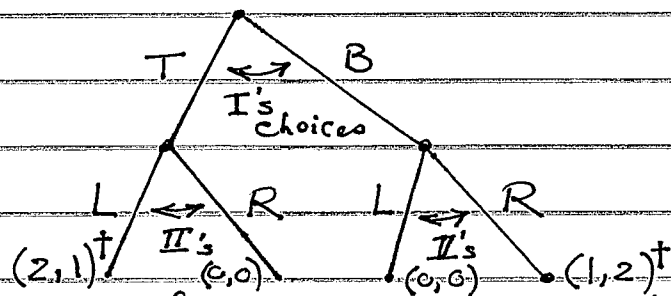
* unstable I \leftarrow II \downarrow
 \uparrow unstable I \uparrow II \rightarrow
^o stable

Battle of Sexes:

		II		
		L	R	
Normal Form	I	T	(2,1) ^o	(0,0)
	B	(0,0)	(1,2) ^o	

^o Stable outcomes (equilibria)

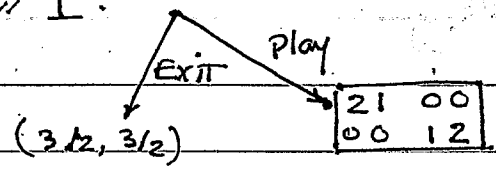
Extensive Form



If this was sequential, i.e. if II chooses after I choose, then II would have chosen \neq options. Assume I can envision this counterfactual mentally. Then I would select T for his choice. This reasoning is known as backward induction.

Exiting in the battle of sexes will yield $3/2$ for I & 2.

Consider the options for I.



Forward induction: I wants to do better than $3/2$, so I plays, hoping for $(2, 1)$ equilibrium.

Hawk-Dove game (HD)

	Hawk	Dove
Hawk	$(\frac{V-C}{2}, \frac{V-C}{2})$	$(V, 0)$
Dove	$(0, V)$	$(\frac{V}{2}, \frac{V}{2})$

ie., a biological game

V = Value of a resource to be enjoyed by a player

Sharing the resource yields $\frac{V}{2}$.

Hawk = tough stance

C = Cost of a conflict with opponent.

if $V > C$, HD becomes the Prisoner's Dilemma, they go for conflict. Cost is relatively low.

if $V < C$, HD becomes the game of "Chicken" (movie Rebel without a Cause)

more generally, this is "Wars of attrition" games. Here fighting is too expensive.